

Class 11: Parsing in Haskell

(guest lecture)

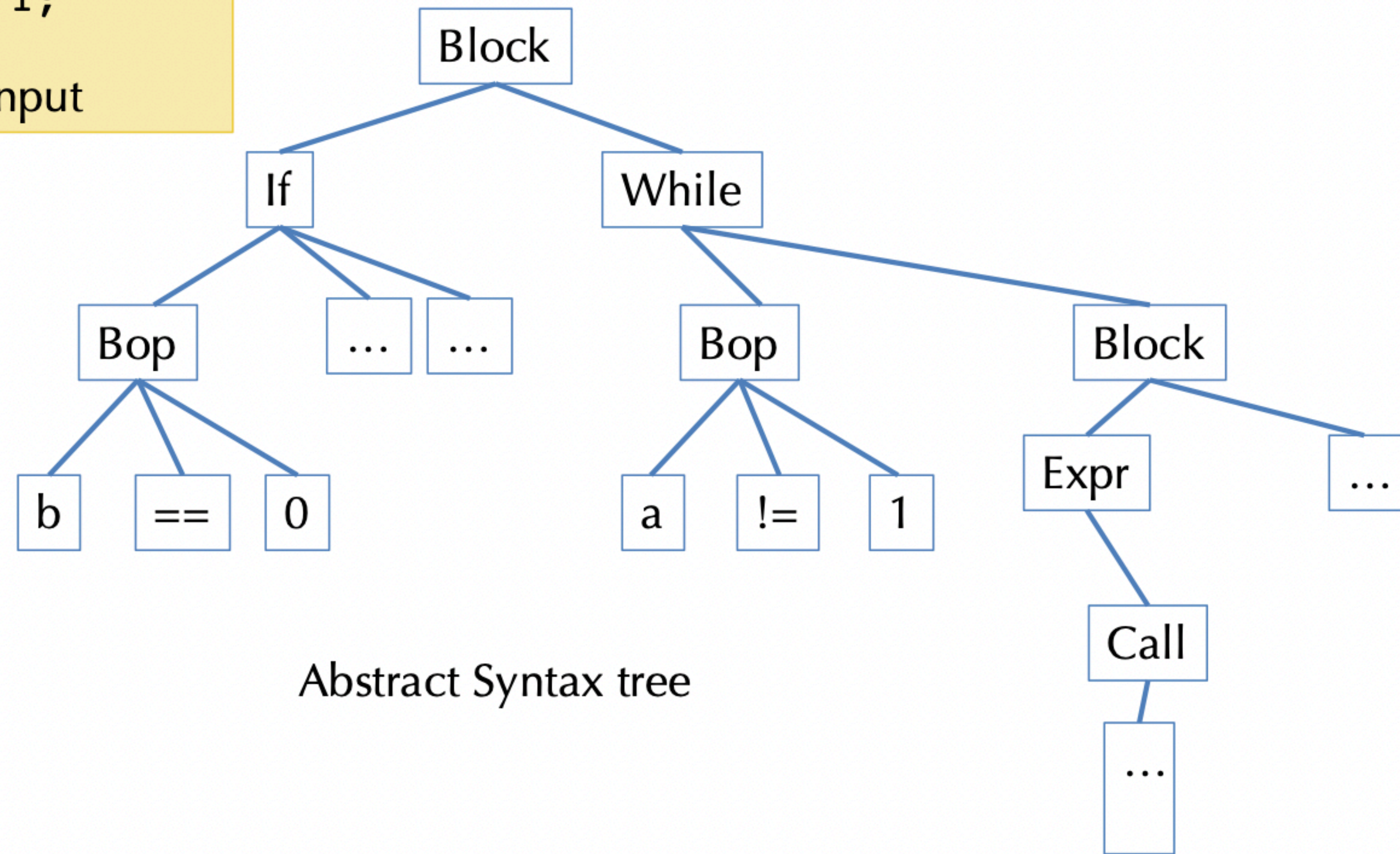
April 4

What is a parser?

type Parser :: String → StructuredObject

```
{  
  if (b == 0) a = b;  
  while (a != 1) {  
    print_int(a);  
    a = a - 1;  
  }  
}
```

Source input



Abstract Syntax tree

Why do parsers matter?

Parsing is a fundamental aspect of many system programming tasks:

CIS 2400 / 3410: Compilers (parse programs)

CIS 3800: Operating Systems (parse terminal input)

CIS 5530: Networks (parsing packets)

Haskell makes writing parsers simple (& fun!)

How do we define a Parser type in Haskell?

Defining the `Parser` type

A Parser is a function:

```
type Parser :: String → StructuredObject
```

A Parser doesn't need to consume all of its input:

```
type Parser = String → (StructuredObject, String)
```

**remainder
of input**

**A Parser should be polymorphic
over the type of structured object that it returns:**

```
type Parser a = String → (a, String)
```

**structured
object**

A Parser should be able to fail (not all strings are parseable!)

```
type Parser a = String → Maybe (a, String)
```

**possibility
of getting
Nothing**

**How do we separate what a Parser *is*
from what a Parser *does*?**

we use *records!*

Crash course on Haskell records

Motivation: how do we give *names* to the arguments of data constructors?

Record example

```
data Student = MkStudent {  
  name :: String,  
  age  :: Int  
}
```

record type definition

**constructor for
the record type**

```
data Student = MkStudent {  
  name :: String,  
  age  :: Int  
}
```

```
ghci> :t MkStudent
```

```
MkStudent :: String → Int → Student
```

name

age

Example: defining a record

MkStudent :: String → Int → Student

```
data Student = MkStudent { name :: String, age :: Int }
```

```
ernest :: Student
```

```
ernest = MkStudent { name = "Ernest", age = 22 }
```

– alternative syntax (order matters!)

```
ernest :: Student
```

```
ernest = MkStudent "Ernest" 22
```

Each record field defines a *selector*

```
ernest :: Student
```

```
ernest = MkStudent { name = "Ernest", age = 22 }
```

```
ghci> name ernest  
"Ernest"
```

```
name :: Student → String
```

```
ghci> age ernest  
22
```

```
age :: Student → Int
```

Defining the `Parser` type

Recall our working type definition for a Parser :

```
type Parser a = String → Maybe (a, String)
```

Let's package this into a record

Separating values of type `Parser a` from their parsing functionality:

```
newtype Parser a =  
  P { doParse :: String → Maybe (a, String) }
```

**actual function that
does the parsing**

```
ghci> :t P  
P :: (String → Maybe (a, String)) → Parser a  
      the doParse function
```

the constructor `P` takes in a function called `doParse` & returns a `Parser`

The **doParse** field in the record defines a record selector:

```
newtype Parser a =  
  P { doParse :: String → Maybe (a, String) }
```

```
ghci> :t doParse
```

```
doParse :: Parser a → (String → Maybe (a, String))
```

The doParse record selector takes a Parser & returns the underlying function

Simple parsers

Parsing one single Char

```
get :: Parser Char
```

```
get = P $ \s →
```

```
  case s of
```

```
    (c : cs) → Just (c, cs)
```

```
    []      → Nothing
```

the stuff after \$ is the doParse function associated with this parser

(in Haskell, you can use \$ to avoid using parentheses)

```
get :: Parser Char
```

```
get = P $ \s → s :: String = [Char]
```

```
case s of
```

```
(c : cs) → Just (c, cs)
```

```
[] → Nothing
```

```
P :: (String → Maybe (Char, String)) → Parser Char
```

Exercise: parsing one single digit

```
oneDigit :: Parser Int
```

Exercise solution: Parsing a single digit

```
oneDigit :: Parser Int
oneDigit = P $ \s →
  case s of
    (c : cs) → do
      i ← readMaybe [c]
      return (i, cs)
    [] → Nothing
```

Exercise solution: Parsing a single digit

```
oneDigit :: Parser Int
```

```
oneDigit = P $ \s →
```

```
  case s of
```

```
    (c : cs) → do
```

```
      i ← readMaybe [c]
```

```
      return (i, cs)
```

```
    [] → Nothing
```

```
s :: String = [Char]
```

**What if instead of `Parsers` for `Ints / Chars`,
we want to build a `Parser` for a function?**

Parsing the sign of integers

```
oneOp :: Parser (Int → Int)
```

```
oneOp = P $ \s →
```

```
  case s of
```

```
    ('-' : cs) → Just (negate, cs)
```

```
    ('+' : cs) → Just (id, cs)
```

```
    _ → Nothing
```

```
oneOp :: Parser (Int → Int)
```

```
oneOp = P $ \s →
```

```
  case s of
```

```
    ('-' : cs) → Just (negate, cs)
```

```
    ('+' : cs) → Just (id, cs)
```

```
    _ → Nothing
```


**What if we want to build a Parser that parses
only when a condition is satisfied?**

Conditional parsing

```
satisfy :: (Char → Bool) → Parser Char
satisfy f = P $ \s → do
  (c, cs) ← doParse get s
  guard (f c)
  return (c, cs)
```

here, we're using the Maybe monad

Parser is a monad

Parser is a Monad

instance Monad Parser where

return :: a → Parser a

return a = P \$ \s → Just (a, s)

(>>=) :: Parser a → (a → Parser b) → Parser b

p >>= k = P \$ \s → do

(a, s') ← doParse p s

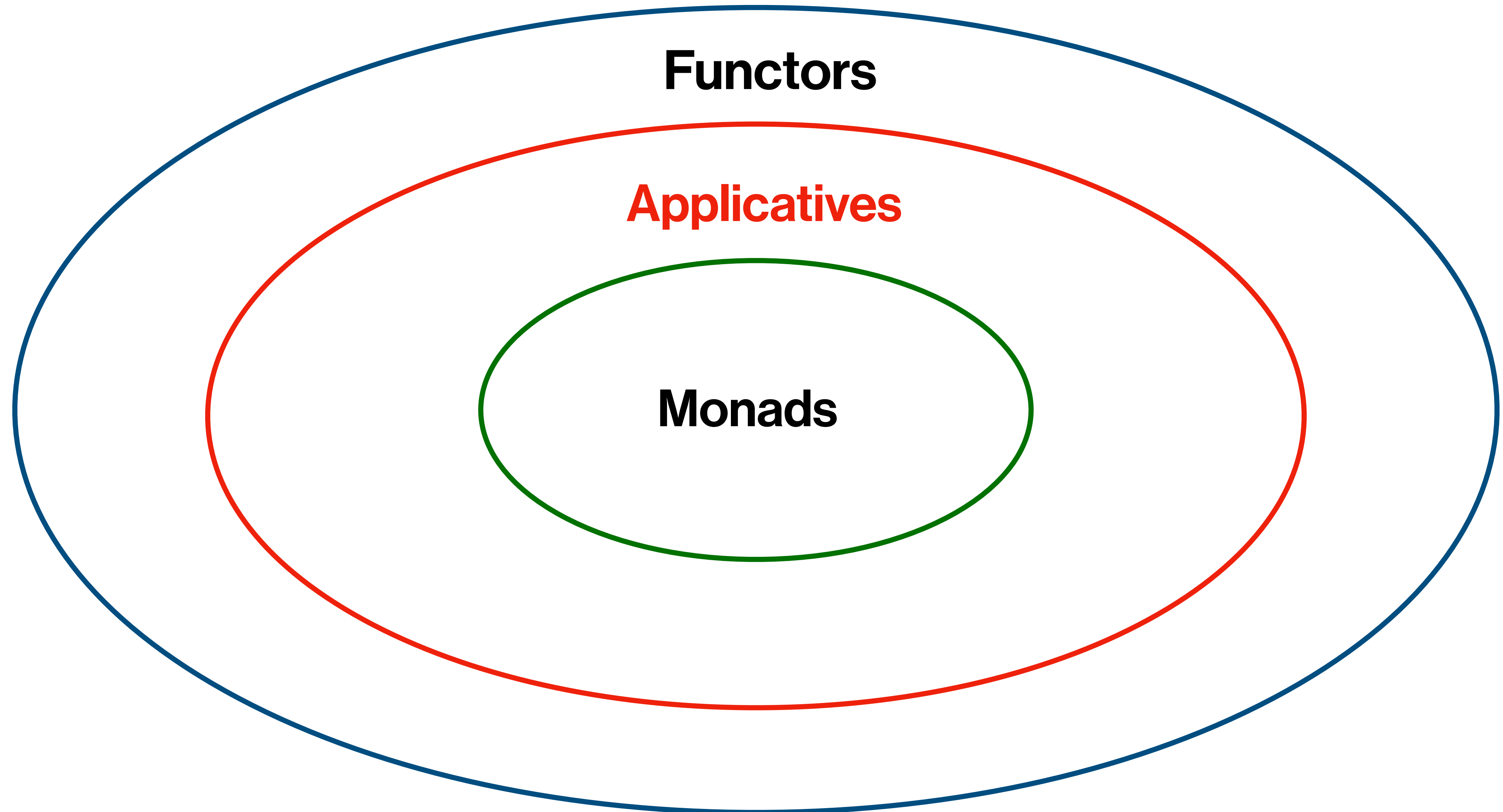
doParse (k a) s'

**As it turns out,
there's a generalisation of Monads that is
more useful for parsing!**

Detour: Applicative functors

(or, how to compose parsers)

Monad \subseteq **Applicative** \subseteq **Functor**



Recall the Functor typeclass:

```
class Functor f where
```

```
  fmap :: (a → b) → f a → f b
```

If you give me:

a “container” of a’s

& a function from a → b

I can give you:

a “container” of b’s

class **Functor** f where

fmap :: (a → b) → f a → f b

What if the **a → b** function itself is contained inside a **Functor**?

Applicative functors
= “functors with application”

```
class Functor f  $\Rightarrow$  Applicative f where  
  pure      :: a  $\rightarrow$  f a  
  (<*)      :: f (a  $\rightarrow$  b)  $\rightarrow$  f a  $\rightarrow$  f b
```

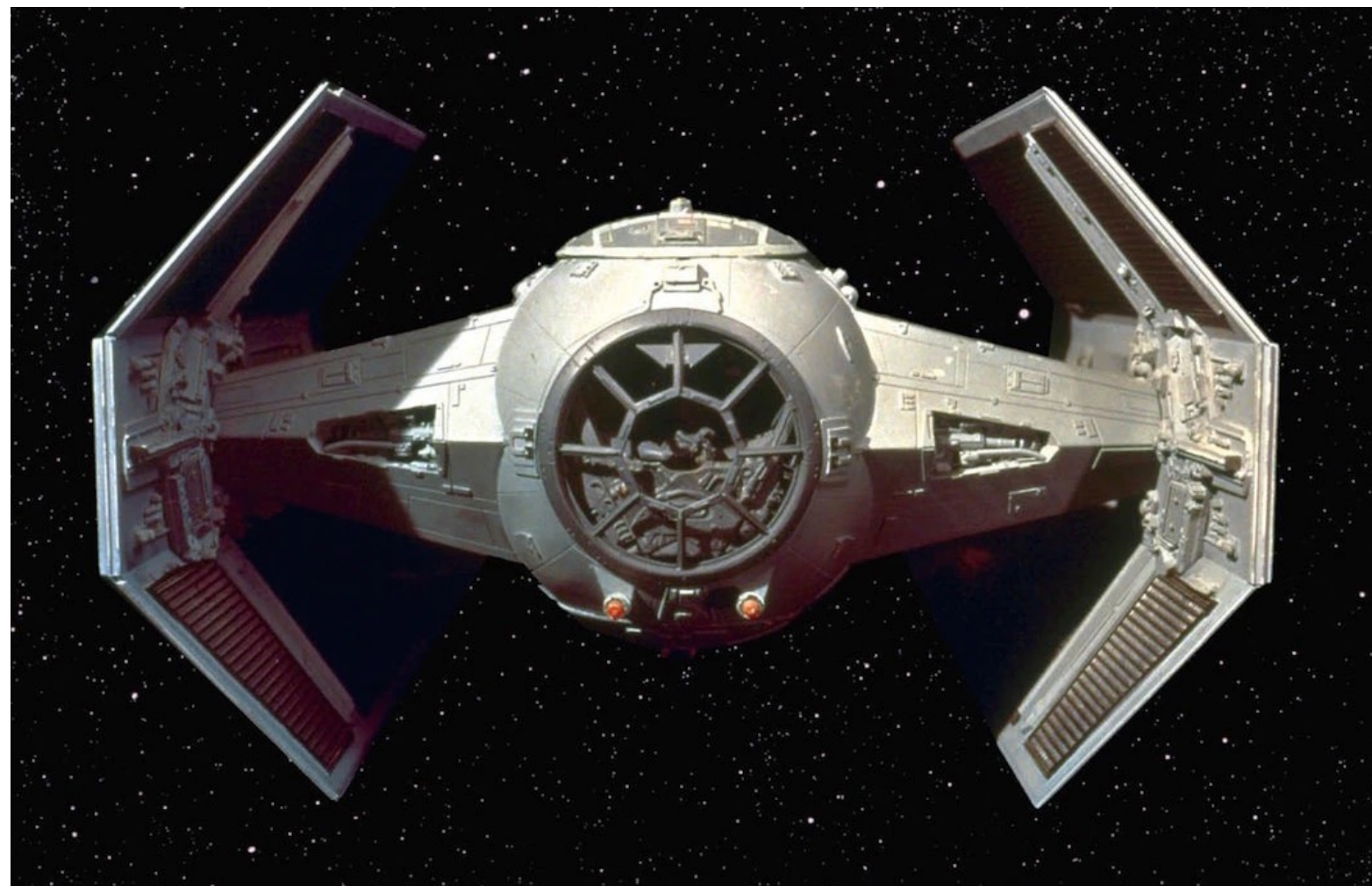
Inject a value into the **Applicative** type

pure :: Applicative f \Rightarrow a \rightarrow f a

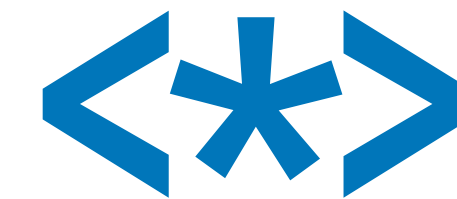
analogous to **return** for Monads

Contextual application of functions

$(\langle * \rangle) :: \text{Applicative } f \Rightarrow f (a \rightarrow b) \rightarrow f a \rightarrow f b$



TIE fighter from Star Wars



“tie fighter”

Parser is an Applicative

```
instance Applicative Parser where
```

```
  pure :: a → Parser a
```

```
  pure a = P $ \s → Just (a, s)
```

```
  (<*>) :: Parser (a → b) → Parser a → Parser b
```

```
  p1 <*> p2 = P $ \s → do  
    (f, s') ← doParse p1 s  
    (x, s'' ) ← doParse p2 s'  
    return (f x, s'' )
```

Composing parsers

$(\langle * \rangle) :: \text{Parser } (a \rightarrow b) \rightarrow \text{Parser } a \rightarrow \text{Parser } b$

$p1 \langle * \rangle p2 = P \$ \backslash s \rightarrow \text{do}$

Composing parsers

$(\langle * \rangle) :: \text{Parser } (a \rightarrow b) \rightarrow \text{Parser } a \rightarrow \text{Parser } b$

$p1 \langle * \rangle p2 = P \$ \backslash s \rightarrow \text{do}$

$(f, s') \leftarrow \text{doParse } p1 \ s$

1. Use parser **p1** to extract a function **f**

Composing parsers

$(\langle * \rangle) :: \text{Parser } (a \rightarrow b) \rightarrow \text{Parser } a \rightarrow \text{Parser } b$

$p1 \langle * \rangle p2 = P \$ \backslash s \rightarrow \text{do}$

(f, s') \leftarrow $\text{doParse } p1 \ s$

(x, s'') \leftarrow **$\text{doParse } p2 \ s'$**

1. Use parser **p1** to extract a function **f**

2. Using the remaining input **s'**, **use parser p2** to extract a value **x**

Composing parsers

```
(<*>) :: Parser (a → b) → Parser a → Parser b  
p1 <*> p2 = P $ \s → do  
  (f, s') ← doParse p1 s  
  (x, s'') ← doParse p2 s'  
  return (f x, s'')
```

1. Use parser **p1** to extract a function **f**
2. Using the remaining input **s'**, use parser **p2** to extract a value **x**
3. **Apply f to x**, and return **s''** (the new remainder of the input)

Example: parsing signed digits

`oneOp :: Parser (Int → Int)`

parses + and -

`oneDigit :: Parser Int`

parses digits

`signedDigit :: Parser Int`

`signedDigit = oneOp <*> oneDigit`

`(<*>) :: Parser (a → b) → Parser a → Parser b`

Monad vs Applicative

Monads vs Applicatives

Monads

$(\gg=)$:: Parser a \rightarrow (a \rightarrow Parser b) \rightarrow Parser b

Applicatives

$(\langle*\rangle)$:: Parser (a \rightarrow b) \rightarrow Parser a \rightarrow Parser b

Monads vs Applicatives

Monads

$(\gg=)$:: Parser a \rightarrow (a \rightarrow Parser b) \rightarrow Parser b

the result of the 1st parser gets to influence how the 2nd parser behaves

Applicatives

$(\langle*\rangle)$:: Parser (a \rightarrow b) \rightarrow Parser a \rightarrow Parser b

the results of the 1st parser are not visible

need to declare beforehand how the parsed objects will be combined (the structure of the computation remains fixed throughout)

Conor McBride, originator of the Applicative typeclass:



"Applicative computations come from the 1960s: you choose all things you're going to compute today and you collect your line-printer output in the morning; **you can't look at the output of one computation and use it to choose a later computation.**"

"Monadic computations come from the 1970s: you sit at your teletype, and **you get to see the response to your previous command before you choose your next command**"

**Monads are more powerful than Applicatives,
but we don't need the full monadic structure for parsing!**

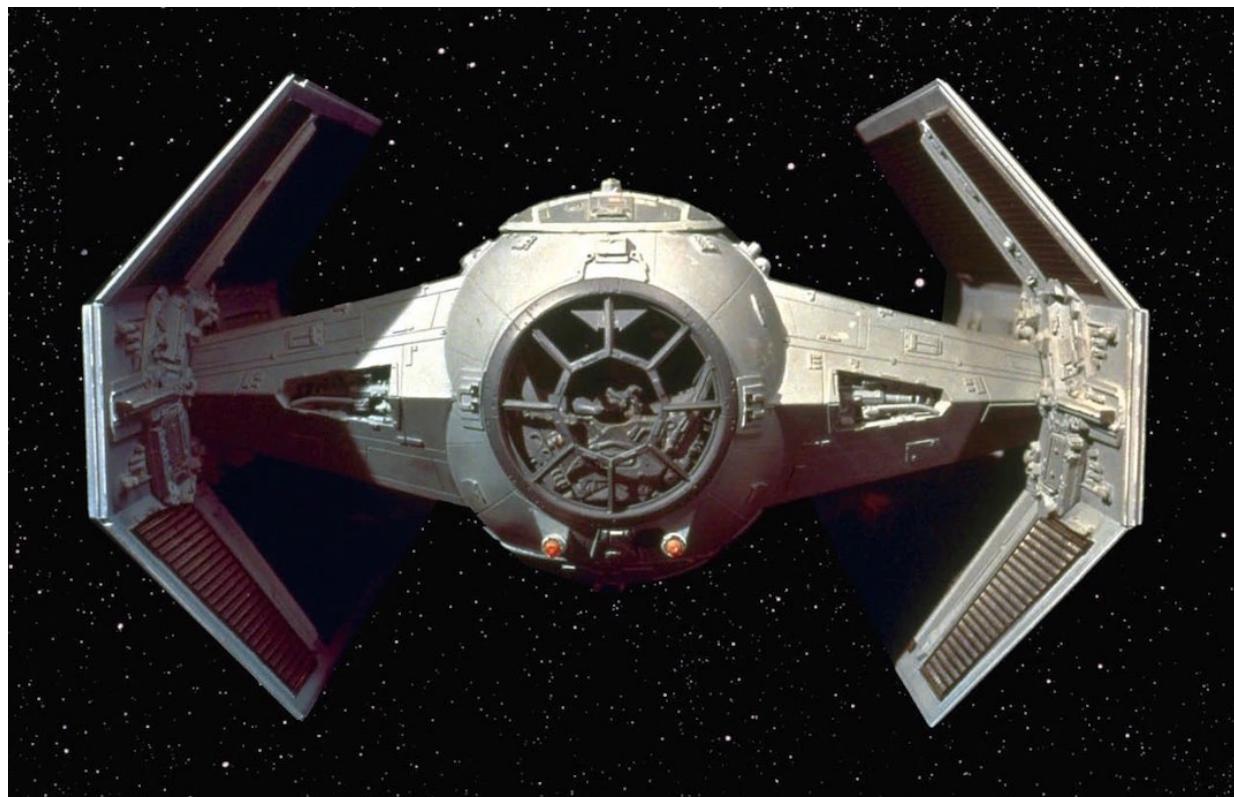
Defining Applicative operators alone is sufficient for parsers*

**(in fact, we only give you an Applicative instance for Parser in this week's HW
— no Monad instance!)**

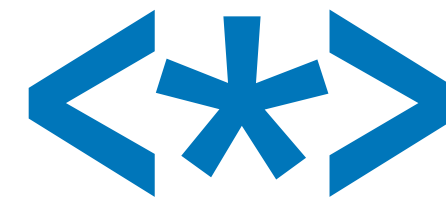
**technically, parsers for any context-free grammar*

Applicative operators

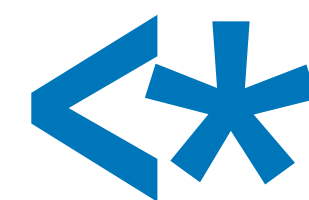
Applicative operators



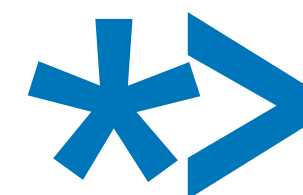
TIE fighter from Star Wars



“tie-fighter”



“left tie”



“right tie”

$(\leftarrow*) :: \text{Applicative } f \Rightarrow \mathbf{f\ a} \rightarrow \text{f\ b} \rightarrow \mathbf{f\ a}$
result is discarded

keep the
1st action's result

$(*\rightarrow) :: \text{Applicative } f \Rightarrow \text{f\ a} \rightarrow \mathbf{f\ b} \rightarrow \mathbf{f\ b}$
result is discarded

keep the
2nd action's result

akin to `>>` for monads

Exercise: using $\ast\rangle$ and $\langle\ast$ to parse between parentheses

`parenP :: Parser a → Parser a`

Example: parsing stuff within (...)

Parser for a specific char

`char :: Char → Parser Char`

`parenP :: Parser a → Parser a`

`parenP p = char '(' *> p <*> char ')'`

parser for (

original
parser

parser for)

Case study:
Translating SQL to Pandas

An automated tool for translating SQL queries into Pandas (joint work with Jason Hom)

<https://github.com/homjason/sql-to-pandas>

(Heavy usage of Applicative-based parsers for parsing SQL)

SQL

```
SELECT day, COUNT(total_bill) as num_bills  
FROM tips_df  
WHERE tip > 2  
GROUP BY day
```



Pandas

```
tips_df[tips_df["tip"] > 2]  
  .groupby(by="day")  
  .agg({"total_bill": "size"})  
  .reset_index()  
  .rename(columns={"total_bill": "num_bills"})
```


Parsing a SQL query

queryP :: Parser Query

queryP = Query <\$> selectExpP

<*> fromExpP

<*> ((Just <\$> whereExpP) <|> pure Nothing)

<*> ((Just <\$> groupByP) <|> pure Nothing)

<*> ((Just <\$> orderByP) <|> pure Nothing)

<*> ((Just <\$> limitP) <|> pure Nothing)

(you'll see in the HW what the (<|>) operator does)

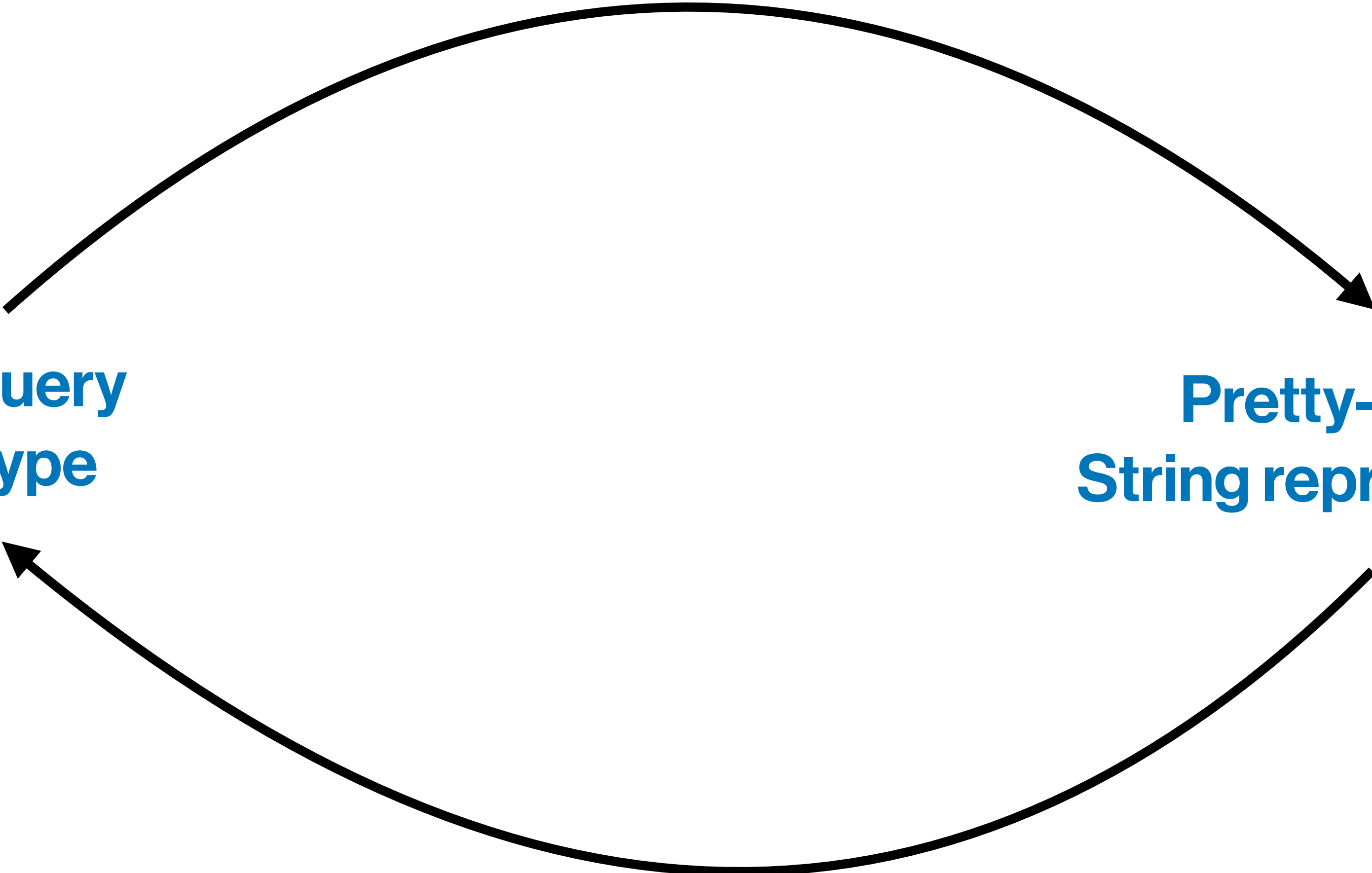
Example QuickCheck round-trip property

Pretty printer

SQL Query
datatype

Pretty-printed
String representation

SQL Parser



Example QuickCheck round-trip property

```
prop_roundtrip :: Query → Bool
prop_roundtrip query =
  doParse queryParser (printQuery query) == Just (query, "")
```

fin

Lecture adapted from:

- **Stephanie Weirich's lectures for CIS 5520**
- **Brent Yorgey's lectures for CIS 1940 (Spring 2013)**